# Workshop on Artificial Intelligence in Practice

Part 1:
Al Targets and Applications in Technics and Logistics

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Section 2a: Algorithms of Dijkstra and A\*

## **Uninformed Search Strategy**

### **Dijkstra's Algorithm for Edge-Valued Graphs**

For all edges (u,v) there is a value: *length* (u,v) := edge evaluation from node u to node v

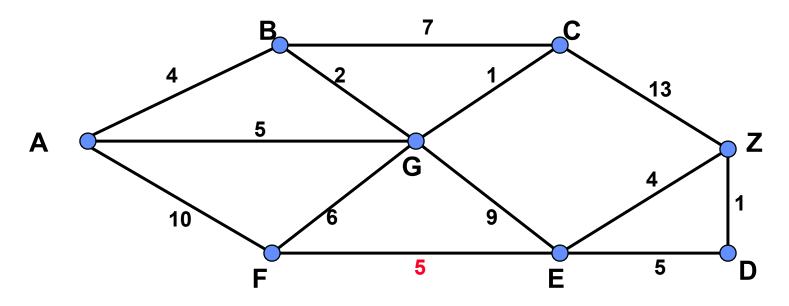
**Prerequesite for edge values:** All edge values must be nonnegative.

#### Algorithm for search of shortest path from A to B (path with minimal sum of edge values):

- Let *Done* be a set initialized by {A}. Set *label* (A) := 0. Let NotFinallyComputed be a set initially consisting of all other nodes. Set *label* (n) := *length* (A,n) for all nodes n directly adjacent to A (the "neighbors"). Set *label* (v) :=  $\infty$  for all nodes n not directly adjacent to A.
- Repeat:

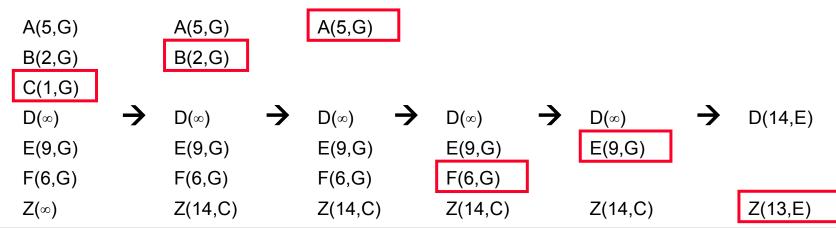
Choose node u from NotFinallyComputed such that label (u) is minimal and move u from *NotFinallyComputed* to *Done*. Update all nodes n from NotFinallyComputed that are directly adjacent to u:  $label(n) := min \{label(n), label(u) + length(u,n)\}.$ until u = B

## **Example for Dijkstra's algorithm**



### **Shortest path from G to Z**: $G \rightarrow E \rightarrow Z$ (13 units)

node (path length from G, direct predeccesor responsible for latest update):



## Informed (heuristic) Search Strategies

### requires the following additional information for each node:

estimation value h (node) as a lower bound for the distance left to the destination

- must be easy to compute (e.g. by geographic coordinates)
- must guarantee that each actual path to the destination is not shorter

h() gives a nonnegative value: The smaller the value, the closer is the distance to the destination

#### Implementation of this principle: A\* Algorithm

- applies Dijkstra's algorithm to this principle
- chooses for shift from NotFinallyComputed to Done not the node with minimal label, but rather the node with minimal sum of label plus estimation value

## Informed (heuristic) Search Strategy

### A\* Algorithm for Edge-Valued Graphs with node heuristic

For all edges (u,v) there is a value: *length* (u,v) := edge evaluation from node u to node v

For all nodes n there is a value:  $h_{\rm B}(n) :=$ lower bound for distance to B

Prerequesite for edge values:

All edge values must be nonnegative.

Prerequesite for node heuristic  $h_{\rm B}(u)$  and actual path length  $p_{\rm B}(u)$  to B:

admissibility:

 $h_{R}(u) \leq p_{R}(u)$ 

(no overestimation)

monotonicity:

 $h_{R}(u) \le h_{R}(v) + length(u,v)$  (triangle inequality)

#### Algorithm for search of shortest path from A to B (path with minimal sum of edge values):

Let *Done* be a set initialized by {A}. Set *label* (A) := 0.

Let *NotFinallyComputed* be a set initially consisting of all other nodes.

Set *label* (n) := *length* (A,n) for all nodes n directly adjacent to A (the "neighbors")

and estimatedPathLength (n) := label (n) +  $h_R(n)$ .

Set label (v) :=  $\infty$  for all nodes n not directly adjacent toand estimated Path Length (v) :=  $\infty$ 

Repeat:

Choose node u from NotFinallyComputed such that estimatedPathLength (u) is minimal and move u from *NotFinallyComputed* to *Done*.

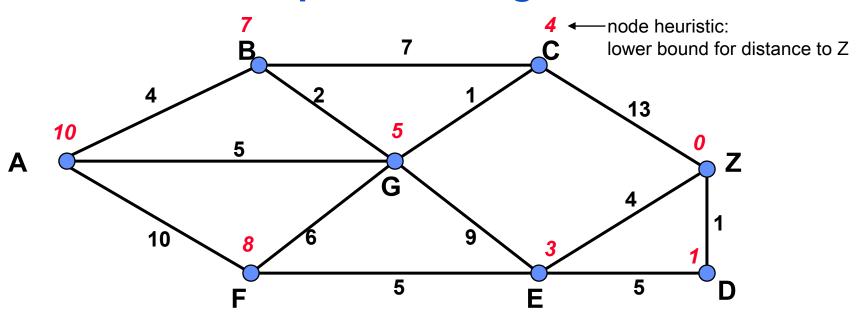
Update all nodes n from *NotFinallyComputed* that are directly adjacent to u:

label (n) := min {label (n), label (u) + length (u,n)}

estimatedPathLength (n) := label (n) +  $h_B(n)$  (if label has changed).

until u = B

## **Example for A\* Algorithm**



### Shortest path from G to Z : $G \rightarrow E \rightarrow Z$ (13 units)

node (path length from G, direct predeccesor responsible for latest update, estimatedPathLength):

