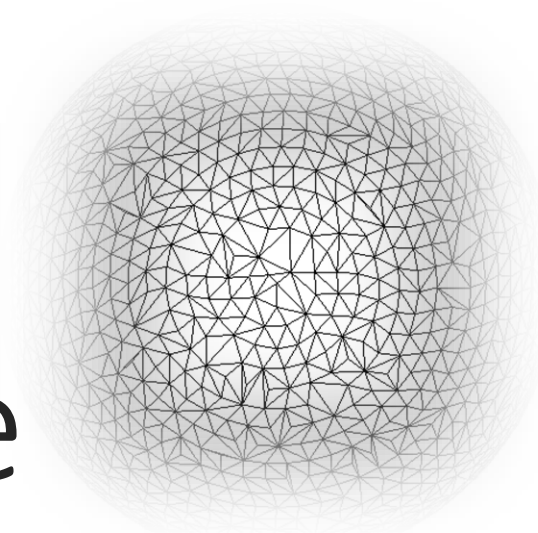


Computational Geometry

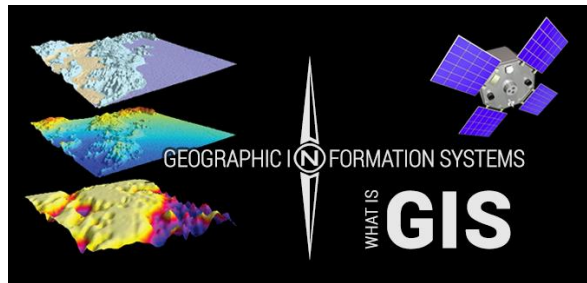
Convex Hull in 3-Space

JONAS SORGENFREI

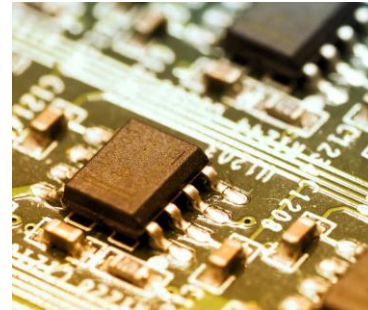


Introduction

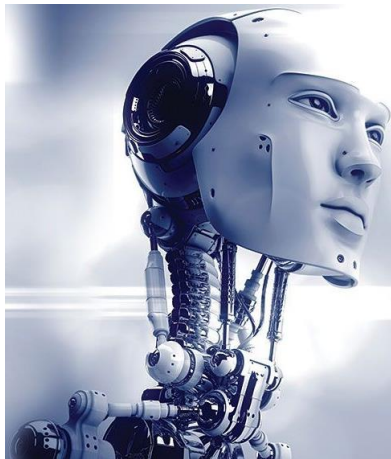
Computational Geometry



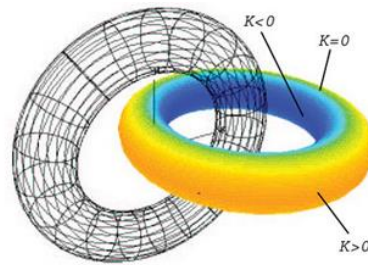
[2] – Geographic Information Systems



[3] – Integrated Circuit design



[1] - Robotics



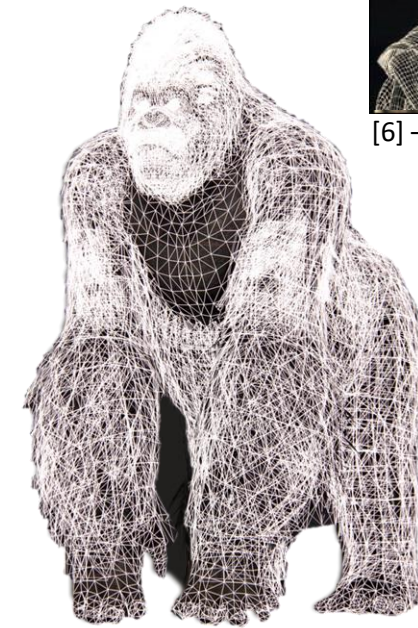
[4] – Computer-Aided Engineering



[5] - Databases



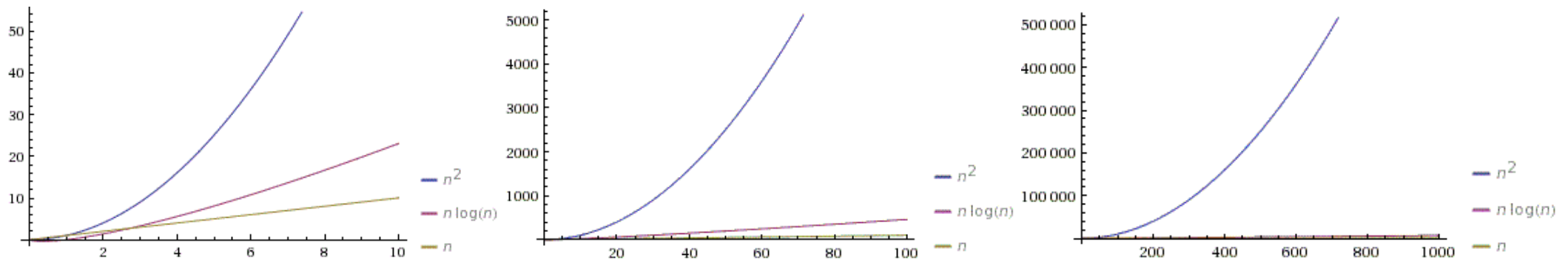
[6] – Computer Vision



[7] – Computer Vision

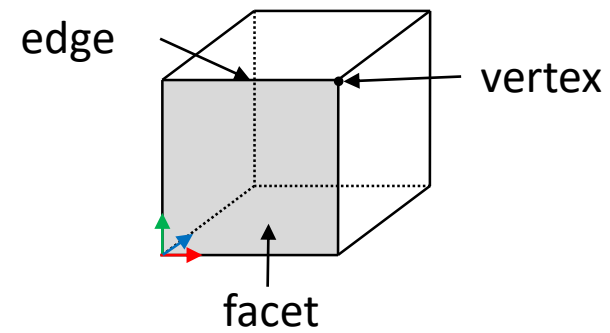
Introduction

Why efficiency matters.



Plots of linear, squared and logarithmic running time

Introduction

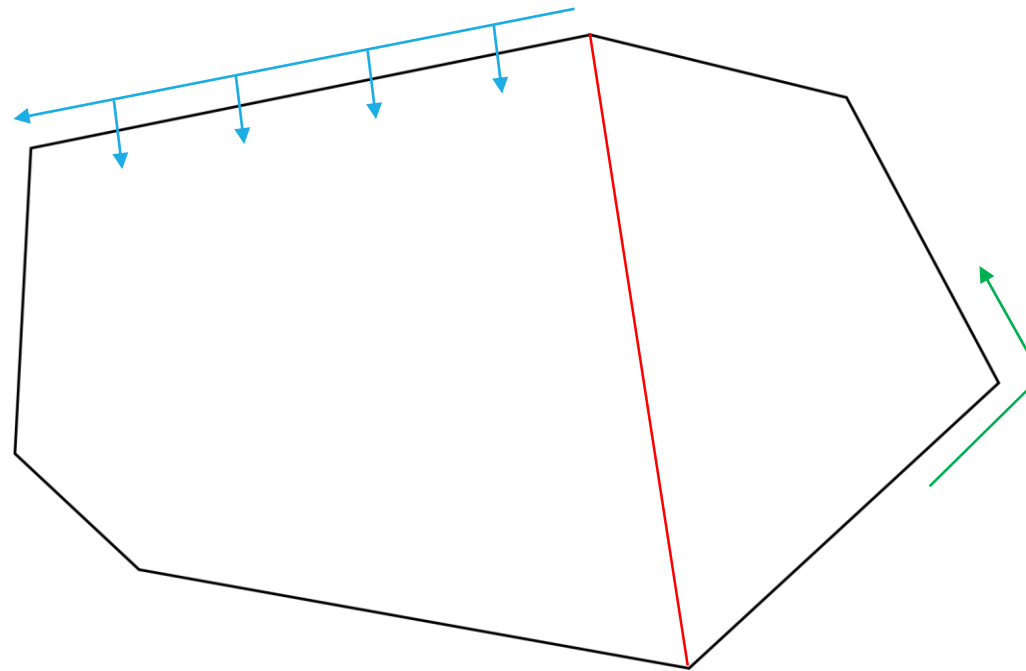


Convex Hull

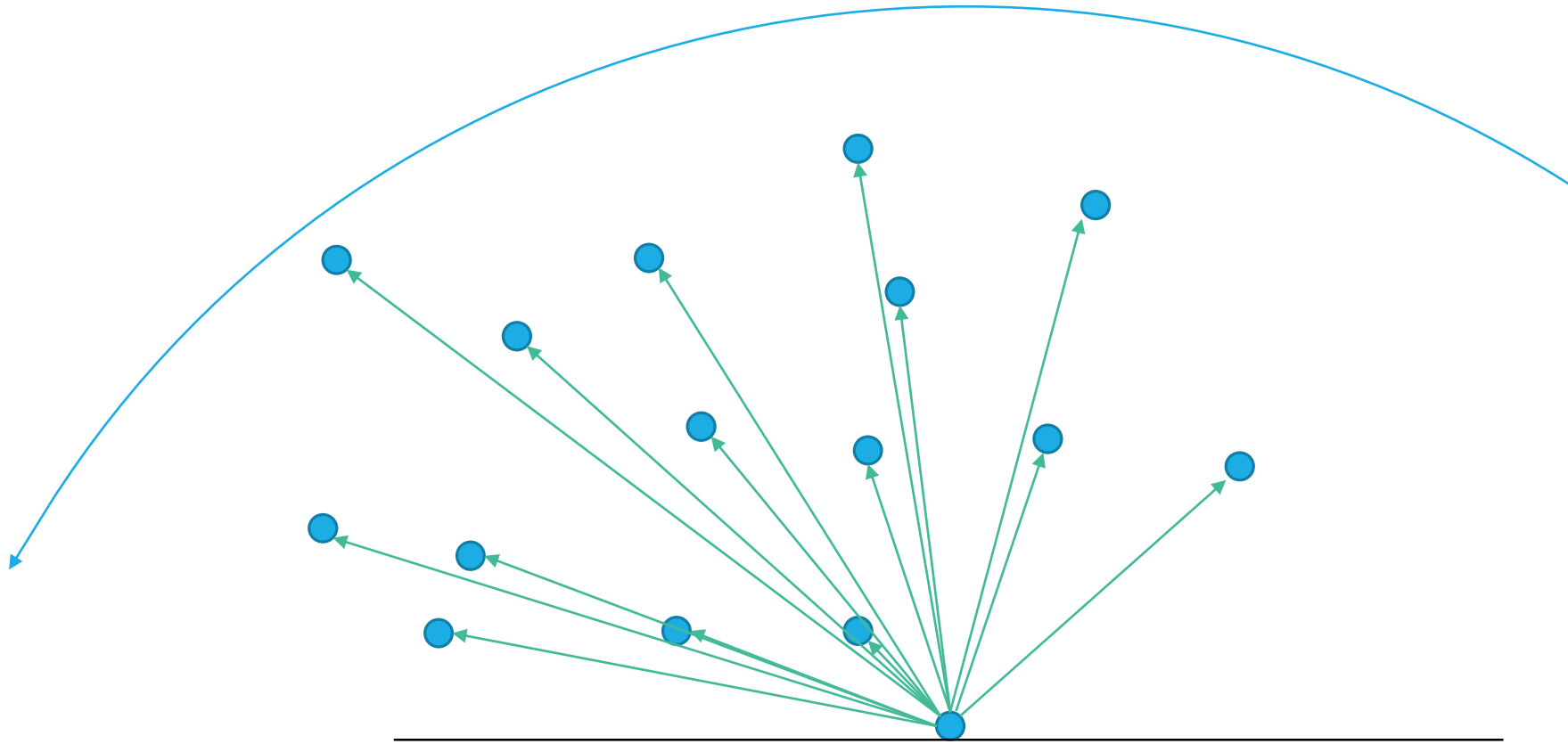
Problem:

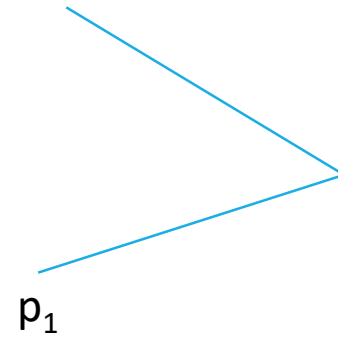
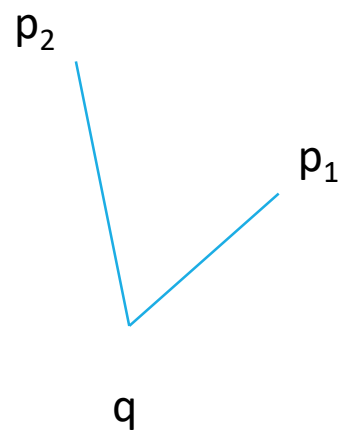
Given a set of n points, find the minimal convex polygon that contains all the points.

Convex Hull - Properties

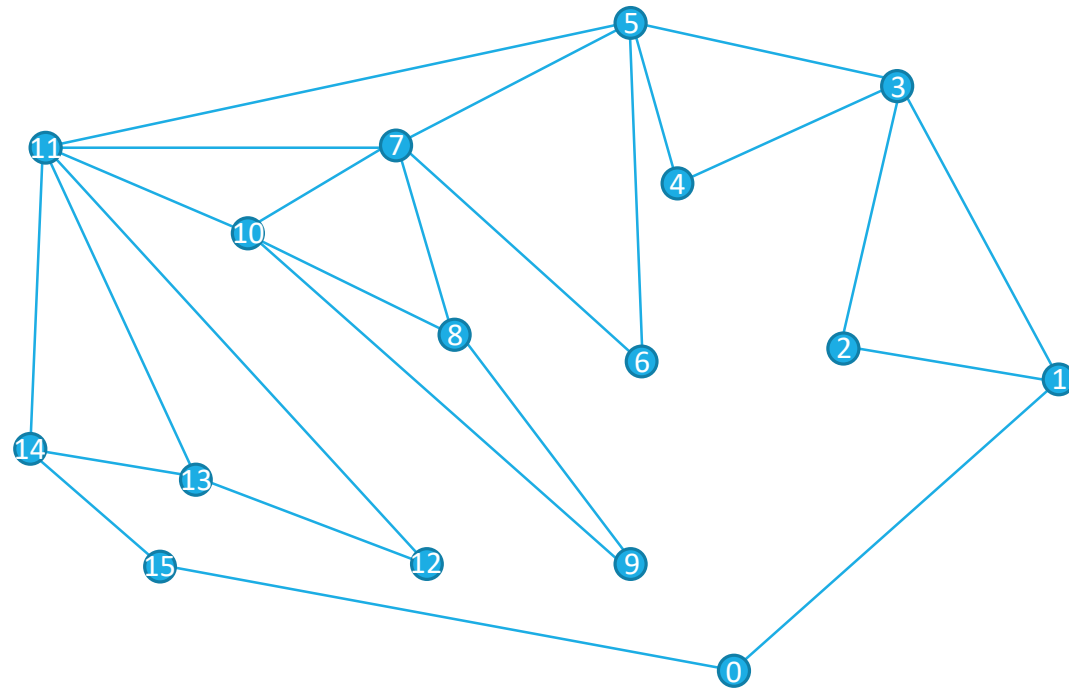


Convex Hull - Algorithms





Convex Hull - Algorithms



Convex-Hull Stack
15
14
5
3
1
0
0

Convex Hull - Algorithms

Pseudo Code – Graham Scan Algorithm

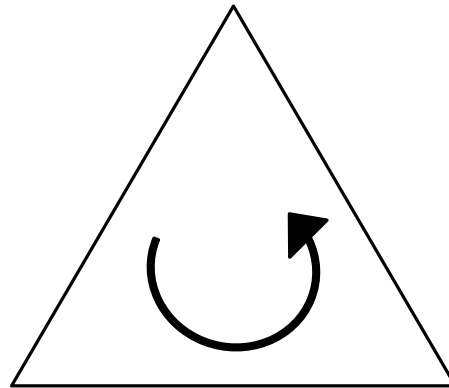
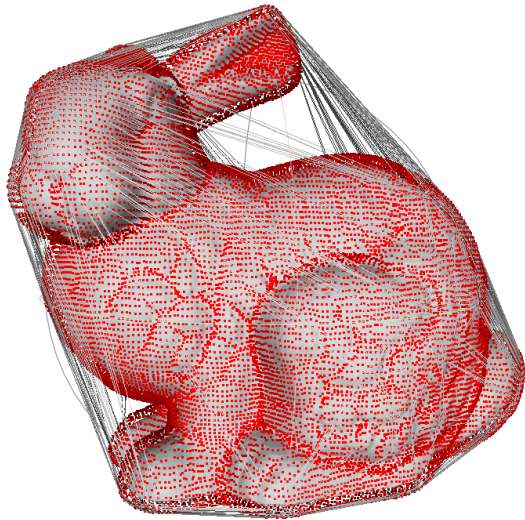
- Select the point with the minimum y and swap with point at [0]
- Sort all points in CCW order $\{p_0, p_1, \dots, p_n\}$
- $S = \{p_0, p_1, p_2\}$
- for $i = 3$ to n
 - While $|S| > 2$ && p_i is to the right of S_{-2}, S_{-1}
 - $S.pop$
 - $S.push(p_i)$

Convex Hull – Algorithms DEMO

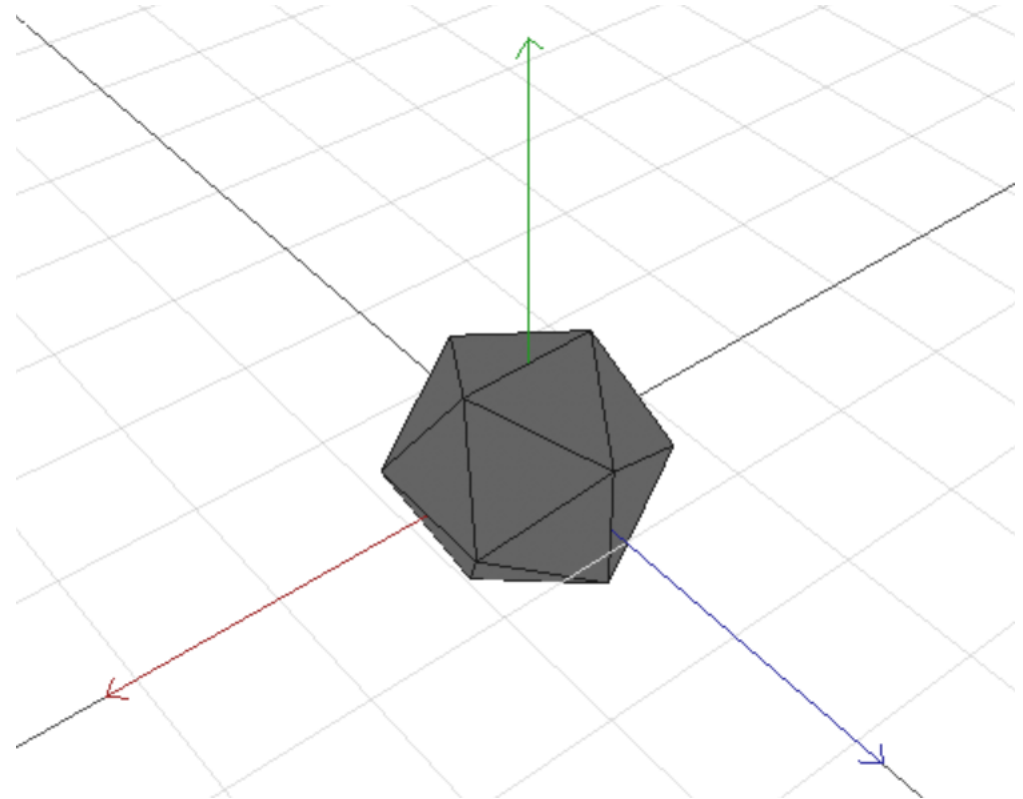


Convex Hull 3D

- polyhedron (convex 3-polytop) containing points
- for every facet of $CH(P)$
 - all other points of P lie in only one half space(back side)

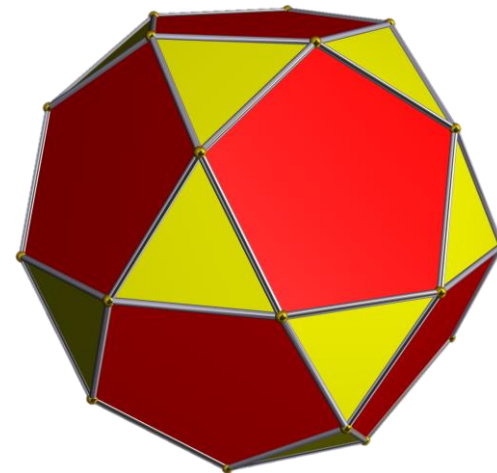
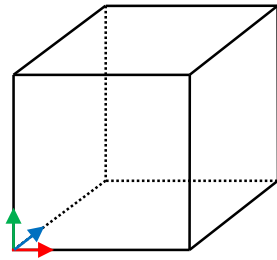


Frontface – CCW-Order



Convex Hull 3D - Complexity

- number of edges can be higher than the number of vertices
- facet = convex polygon



Icosidodecahedron

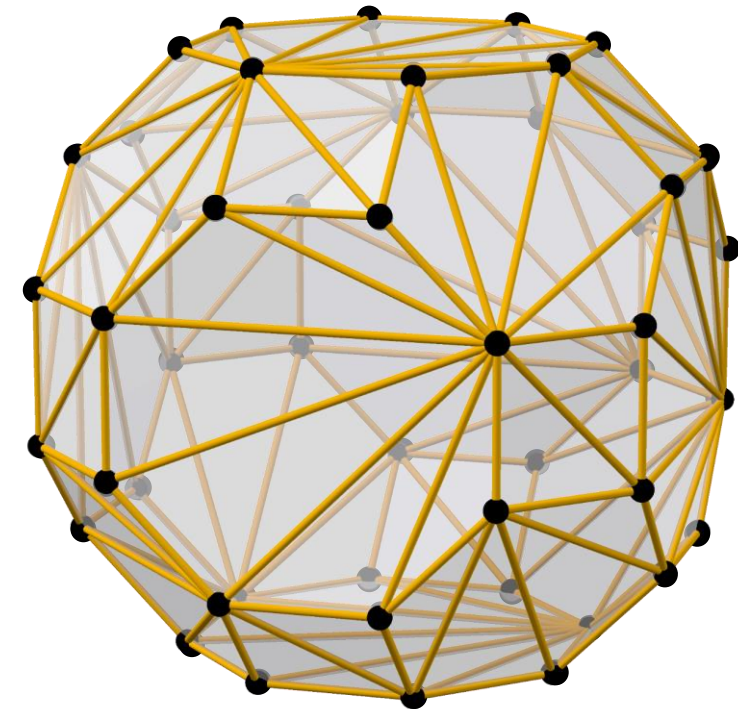
Convex Hull 3D - Computation

2 main problems

- finding extreme points
- finding connectivity

Saving CH in a Data Structure

- Saving Points
- Saving Facets (Connection of points)



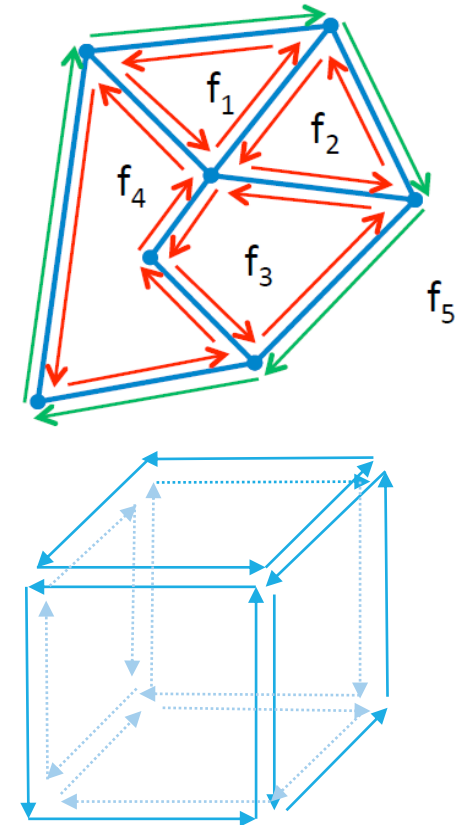
Doubly Connected Edge List

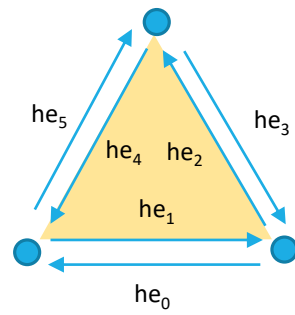
- data structure
- represents an embedding of a planar graph in the plane and polytopes in 3D
- needs functions to work on abstract data structure

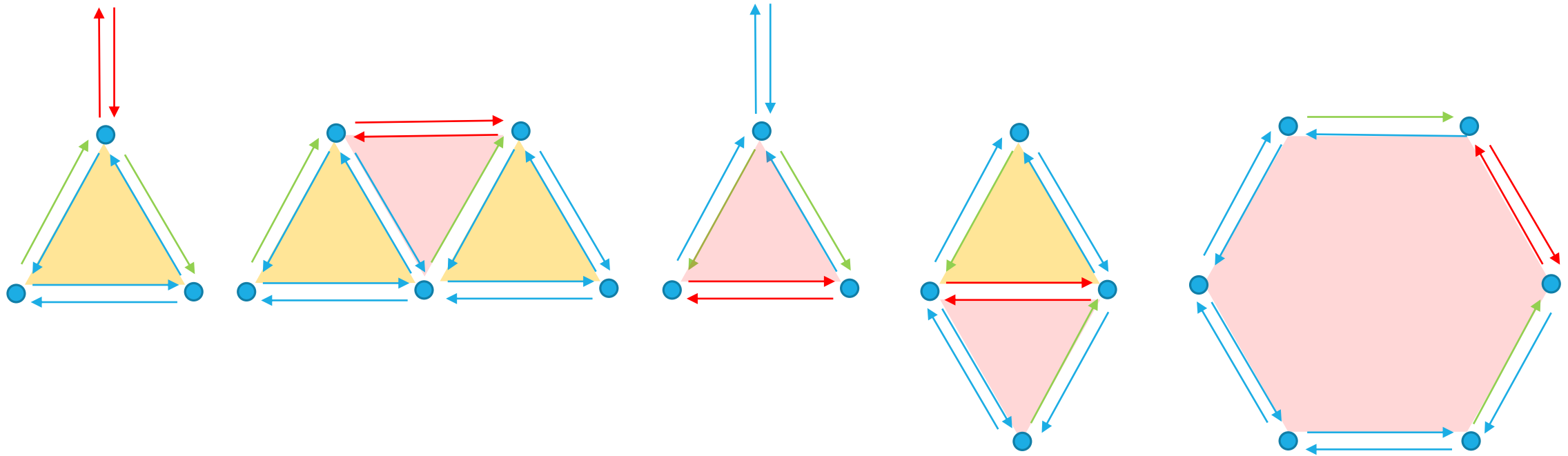
links together:

- **Vertex**
- **Halfedge**
 - Edge Connecting $P1 \rightarrow P2$
 - (Twin) Edge Connecting $P2 \rightarrow P1$
- **Face**

complex & complicated structure/programming!







Convex Hull 3D - Computation

Randomized incremental algorithm

Input

$P = \text{Set}(\text{point}_1, \text{point}_2, \dots, \text{point}_n)$

Output

$\text{CH}(P)$ (*Convex Hull of P*)

Convex Hull 3D - Computation

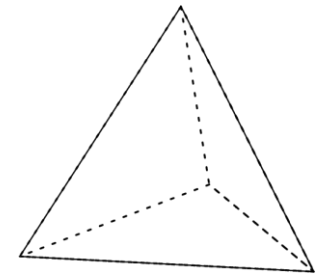
1) choosing 4 Points, that **don't** lie in a common plane (tetrahedron)

- p1 and p2 be 2 Points in P
- p3 a point in P that does not lie on the line of p1 and p2
- p4 a point in P that does not lie in the plane p1,p2,p3
- (can't find 4 points \rightarrow all point in P in a plane \rightarrow 2D CH algorithm)

2) compute random permutation p_5, \dots, p_n of remaining points

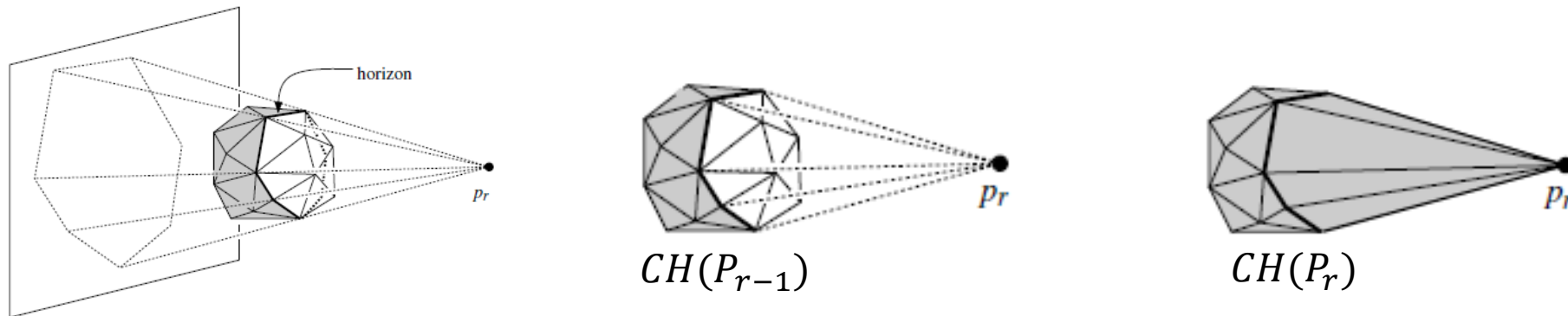
3) for all points p_5, \dots, p_n

- In a generic step of the algorithm we have to add the point p_r to the convex hull of $CH(P_{r-1})$
- \rightarrow transform $CH(P_{r-1})$ to $CH(P_r)$
 - \rightarrow if p_r in $CH(P_{r-1})$ or on it's boundary $\Rightarrow CH(P_r) = CH(P_{r-1})$
 - \rightarrow else p_r outside of $CH(P_{r-1})$ adding new point deleting old ones



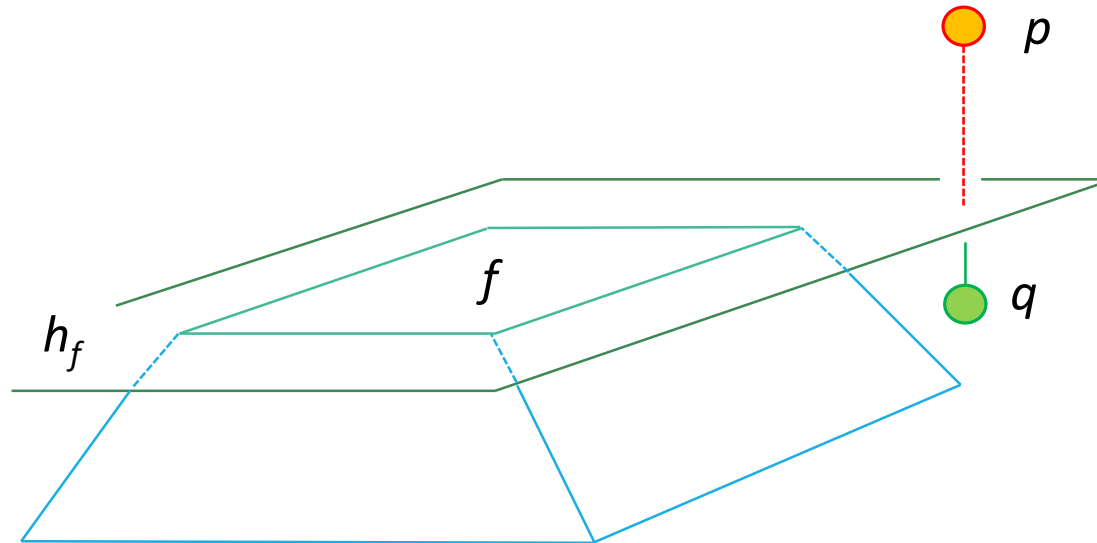
Tetrahedron

Convex Hull 3D - Computation



Convex Hull 3D - Computation

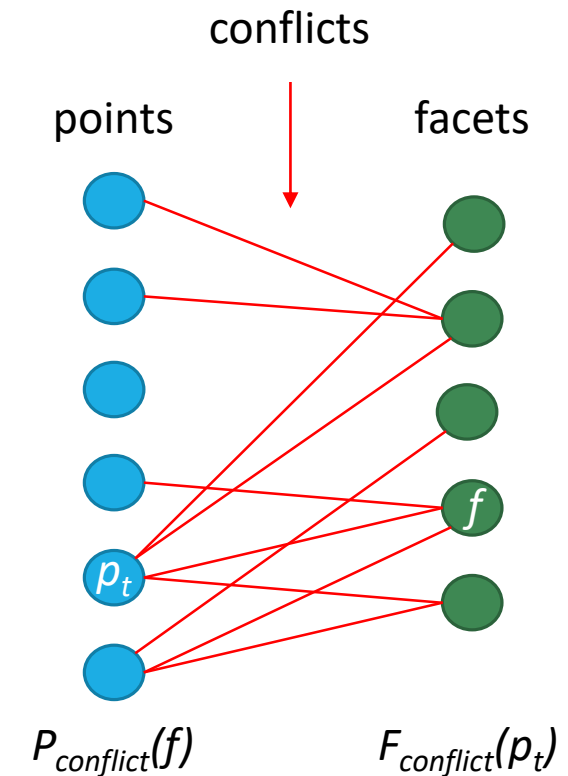
f is visible from p ,
but not from q



Conflict Graph

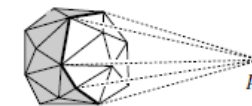
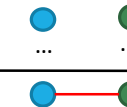
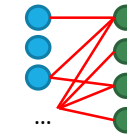
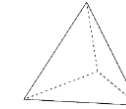
Data structure to store

- visible facets $\{f_{n'}, f_{m'} \dots\}$ from point p_t
 - as conflicts
- visible points $\{p_{n'}, p_{m'} \dots\}$ from point f_t
 - as conflicts
- methods for finding Conflicts-Partners
 - $P_{\text{conflict}}(f)$ – List of Points in Conflict with Face f
 - $F_{\text{conflict}}(p_t)$ – List of Faces in Conflict with Point p_t



Convex Hull 3D – Algorithmen Overview

- Starting with a Tetrahedron
 - 4 non-Coplanar points of the Set P
- Compute random permutation for remaining points
- Initialize the Conflict-Graph with Tetrahedron & remaining points
- Loop over all remaining points
 - Check if Conflict Graph of current point is Empty
 - True: continue with next point (point is inside current C. Hull)
 - False:
 - Delete Faces & Build up the Horizon
 - Go along the Horizon & Create new Faces with the current Point
 - Check for new Conflicts between new Face & remaining Points (not c. Point)
 - Delete Old Conflicts

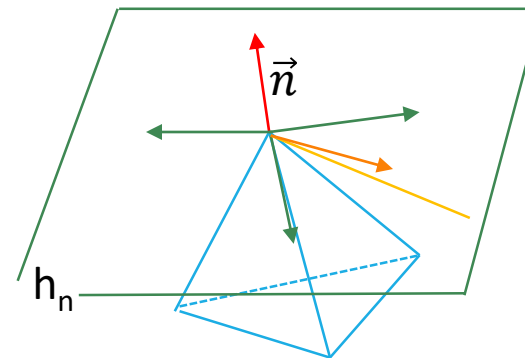


Sample Program

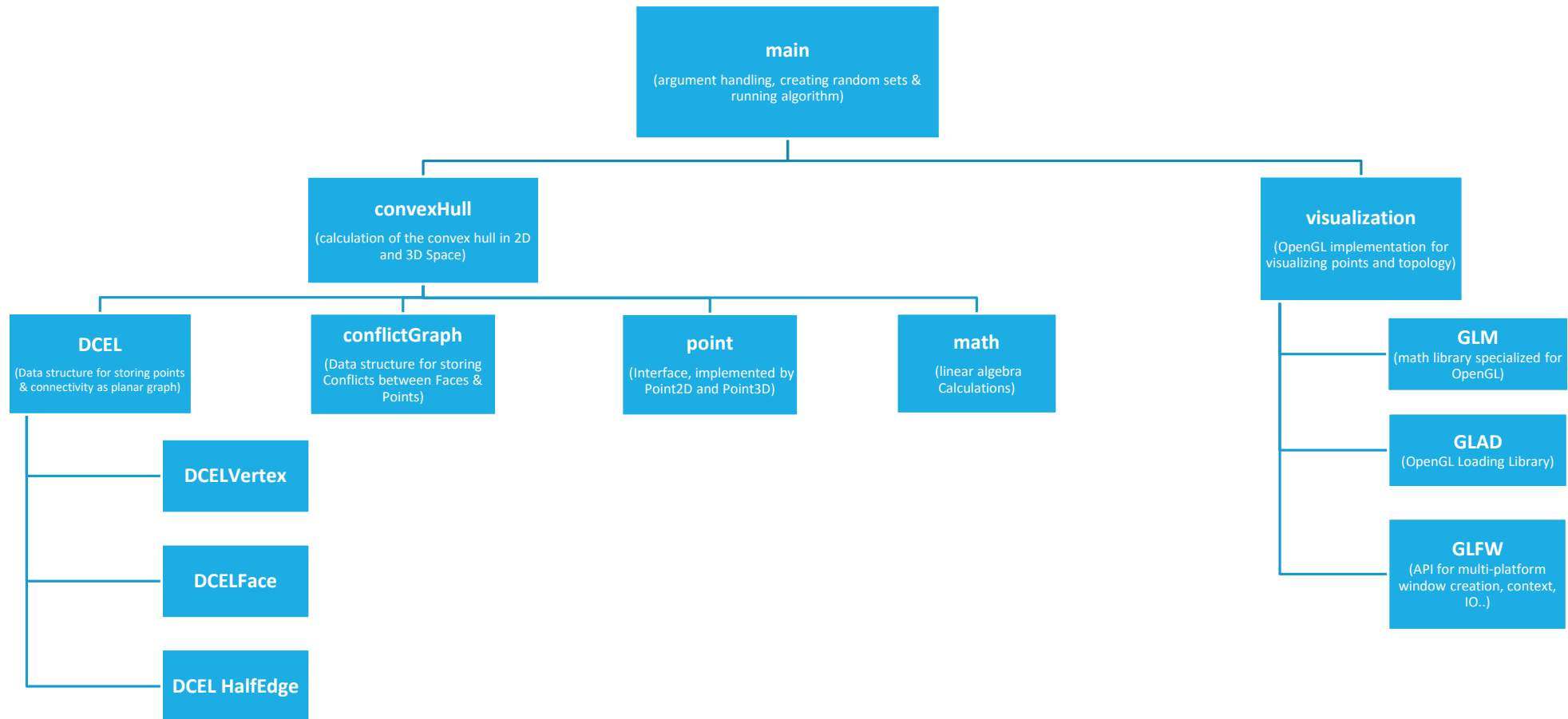
Implementation of a 3D convex hull program

- 2D & 3D convex hull
- parameter arguments as runtime-arguments (-h → help shows all possible arguments)
- using OpenGL for the visualization

<https://github.com/jonassorgenfrei/convexHull3D>



Sample Program



Convex Hull 3D - Conclusion

- running time (expected) $O(n \log n)$
- generalizes to **higher dimensions** (optimal in the worst case)
- best **deterministic** convex hull algorithm for odd-dimensional spaces is based on a (quite complicated) de-randomization of this algorithm
- improvements for applications
 - parallelize algorithm → use on Hardware (GPU)



Computer Graphic Applications

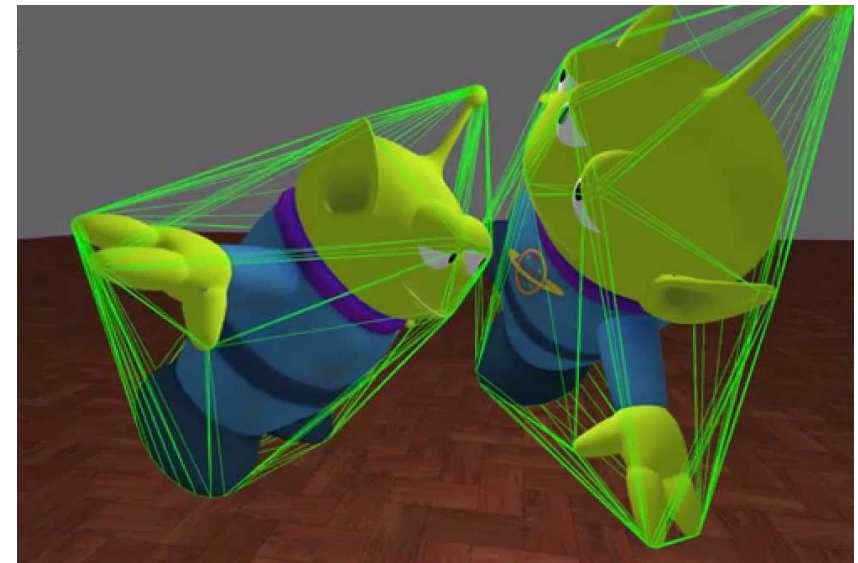
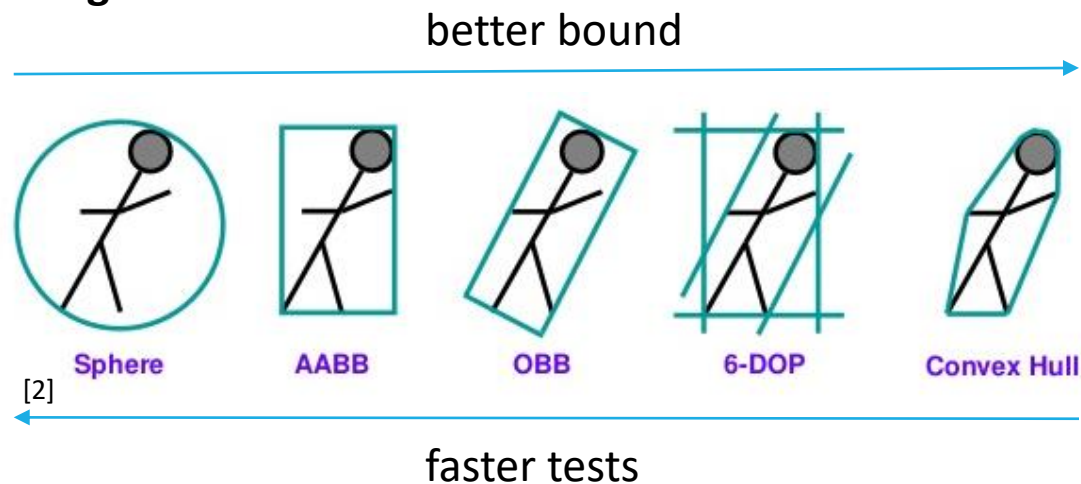
Convex Hull 3D

Computer Graphic Applications

Collision Detection in Computer Animation / Computer Games

- speed up (when result is negative most of the time)
- less costly than actual polyhedron
- better approximation

Bounding Volumes



[1]

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Convex Hull in 3D Dimensions – Peter Felkel Fel CTU Prague – Version 23.10.2014

GPU accelerated Convex Hull Computation - Min Tanga, Jie-yi Zhaoa, Ruo-feng Tonga, Dinesh Manochab

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Additional Ref

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External Graphics

Slide 3

- [1] <https://www.techiexpert.com/future-proof-clients-portfolios-robotics-ai/>
- [2] <https://gisgeography.com/what-gis-geographic-information-systems/>
- [3] <http://www.imperial.ac.uk/continuing-professional-development/short-courses/eng/electrical/cmos/>
- [4] <https://ocw.mit.edu/courses/mechanical-engineering/2-158j-computational-geometry-spring-2003/>
- [5] <https://www.cbronline.com/what-is/what-is-a-database-4917209/>
- [6] <https://techterms.com/definition/rendering>
- [7] <http://www.zbrushcentral.com/printthread.php?t=168180&pp=15&page=16>

Slide 13

<http://diskhkme.blogspot.com/2015/10/convex-hull-algorithm-in-unity-2-3d.html>

Slide 14

<https://upload.wikimedia.org/wikipedia/commons/0/02/Icosidodecahedron.png>

Slide 17

https://commons.wikimedia.org/wiki/File:Truncated_cuboctahedron,_ball-and-stick,_triangles.png

External Graphics

Slide 33

<http://www.cgrecond.net/2016/02/deadpool-vfx-breakdown.html>

Slide 35

[1] <https://www.youtube.com/watch?v=O4FC5rxEqGk>

[2] <https://www.slideshare.net/oiotext/collision-detection-ooxv0p2>

Slide 36

<https://www.sidefx.com/company/press/>